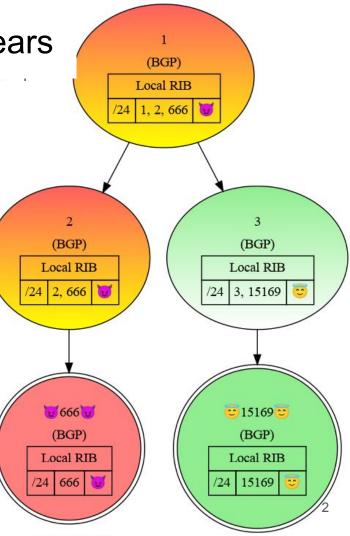
Securing BGP ASAP: ASPA and Other Post-ROV Defenses

By: Dr. Justin Furuness, Dr. Cameron Morris, Dr. Reynaldo Morillo, Arvind Kasilya, Dr. Bing Wang, Dr. Amir Herzberg

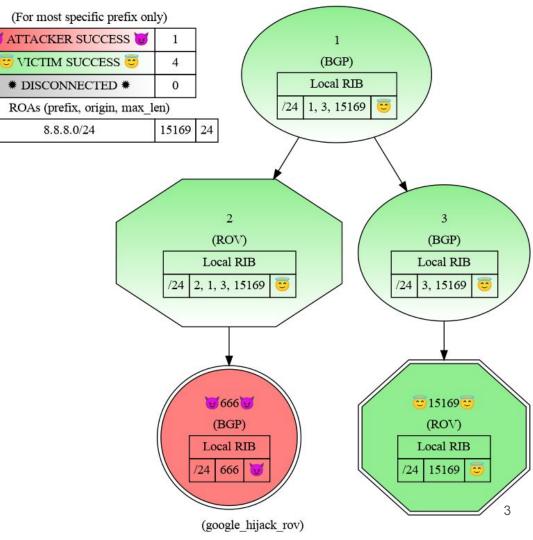
BGP Hijacks: A Problem for over 30 years

- The internet is comprised of over 100k Autonomous
 Systems (ASes) and >500k connections
- These ASes communicate using the Border Gateway Protocol (BGP)
- From a high level, ASes announce which prefixes they own to the rest of the internet
 - For example, Google's AS of 15169 announces prefixes such as 8.8.8.0/24
- BGP was not designed with security in mind
 - Other ASes can claim to own google.com and reroute the traffic through themselves
 - These hijacks happen nearly every day: <u>https://x.com/bgpstream</u>



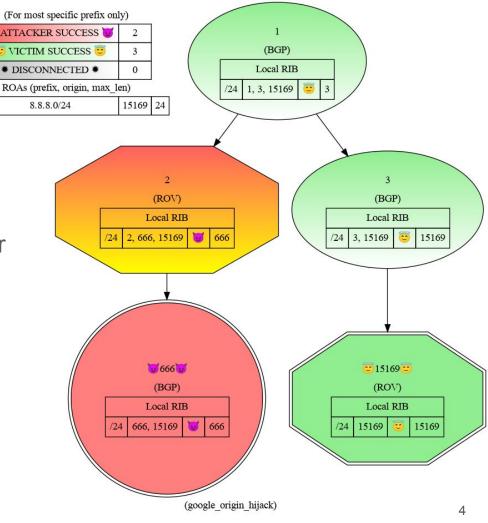
Route Origin Validation (ROV)

- Internet Engineering Task Force (IETF) proposed standard to secure BGP
- From a high level, you can think of ROV using the RPKI to keep a list of valid prefix-origin pairs
 - The origin here is the AS number (or ASN for short)



Forged-Origin Hijacks

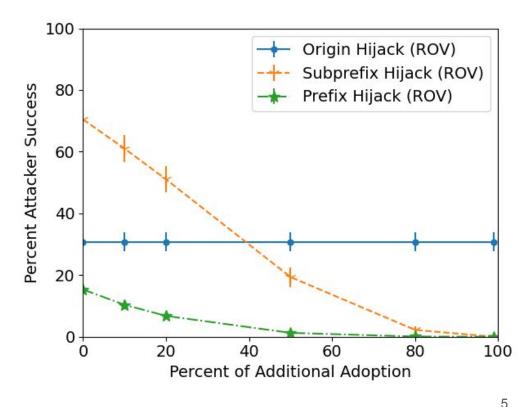
- While ROV deployment is widespread, unfortunately, ROV does not cover all types of attacks
- In a forged-origin hijack, the attacker can fake the AS-Path, and bypass ROV
- The attacker does incur the penalty of a longer AS-Path
- When sent to the attacker's providers, this powerful hijack propagates to the entire internet



Local RIB rows displayed as: prefix, as path, origin, next_hop

Is ROV Good Enough?

- To test this, in a first, we simulated the current internet with real-world AS and ROV data
- We discovered that already, today, forged-origin hijacks are more powerful than prefix-hijacks, even with ROV deployment!
 - Subprefix hijacks (which we don't have time to get into) are not possible for /24 prefixes
- So how can we defend against this powerful forged-origin attack?

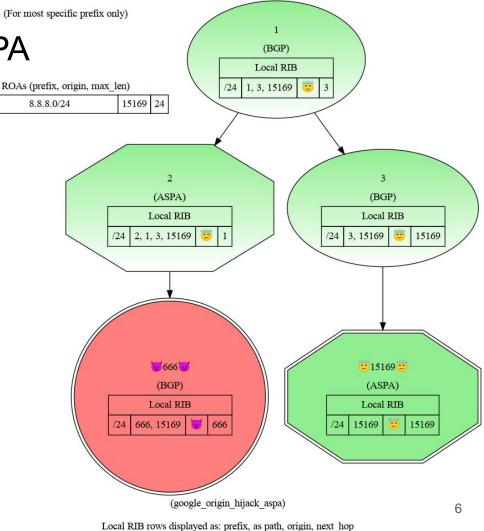


Categories 3, 6, 7 from "Keep Your Friends Close, but Your Route Servers Closer"

(For most specific prefix only)

IETF Proposed Defense: ASPA

- ASPA functions by publishing a list of your providers in ASPA records
 - Customer-provider Ο
 - Ο peer-peer
- Whenever an ASPA AS receives a BGP announcement from a customer or peer:
 - If the provider chain is not Ο possible, the announcement is rejected
 - ASPA also protects against leaks Ο
- ASPA is readily deployable and already in the wild!

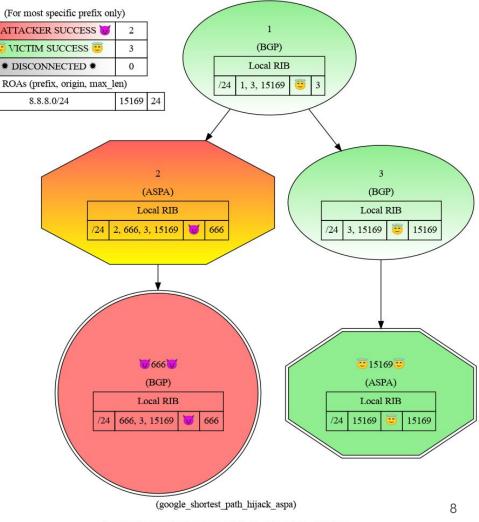


Securing BGP ASAP: ASPA and Other Post-ROV Defenses

- In our work, we evaluate ASPA against other Post-ROV defenses (that can be found in our paper) and propose extensions to it that we call ASPA with Neighbors (ASPAwN)
- We find that ASPA is very effective at preventing:
 - Forged-Origin hijacks
 - Shortest-Path hijacks (we are the first to evaluate this)
 - Route Leaks
 - In a first, we find that ASPA is better than OTC (another IETF RFC)!
- Prior works came to the conclusion that ASPA is not effective at intermediate ASes, or without tier-1 adoption a conclusion that may deter would-be adopters.
 - In our work, we show that ASPA is effective regardless of tier-1 adoption, motivating ASPA adoption at all transit ASes
- We discuss vulnerabilities relating to the novel First-ASN-Stripping Hijack
- Our work is fully reproducible and open-source (links at the end)

Shortest Path Hijack

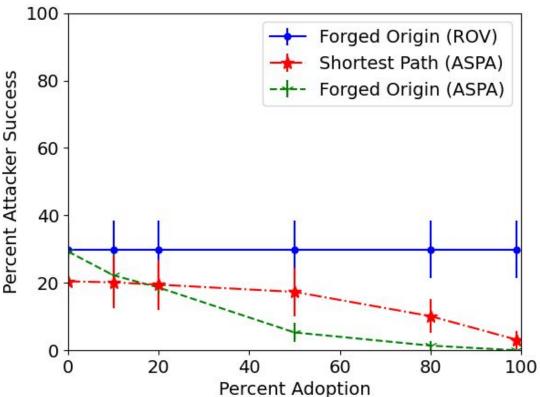
- There is another type of hijack designed to work against ASPA
- The Shortest-Path hijack keeps adding ASes to the AS-Path until there is an AS that does not adopt ASPA
 - When this occurs, ASPA can not detect the hijack, since the last AS in the AS-Path does not announce their providers and does not have ASPA records



Local RIB rows displayed as: prefix, as path, origin, next_hop

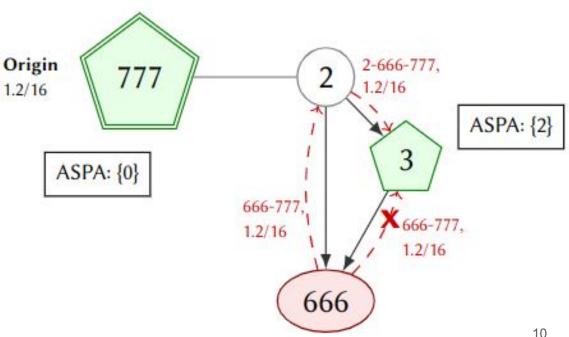
Forged-Origin and Shortest-Path Hijack Simulations

- As we can see, ASPA is highly effective against both of these attacks
 - Significantly better than
 ROV
- Initially the forged-origin hijack is the strongest due to the shorter AS-Path
- Once ASPA adoption is high, the shortest-path hijack is stronger
 - (Since this bypasses ASPA with the longer AS-Path)



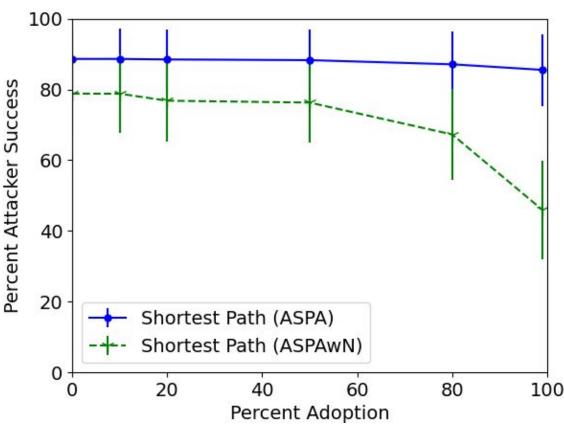
ASPA with Neighbors (ASPAwN)

- ASPA only validates announcements sent to providers
- ASPA rejects the hijack from a customer - but accepts from a 1.2/16provider!
- ASPAwN adds an optional set of valid next-hops to the ASPA record to prevent this
- This was developed in parallel to ASRA and has equivalent security properties



Shortest-Path Hijack from a Transit AS

- Here a transit AS performs the shortest-path hijack against ASPA and ASPAwN adopting ASes
- We are measuring the customer cone of the transit AS
- ASPAwN is significantly more effective at protecting customers
- (About 50% of ASes go through the attacker passively)



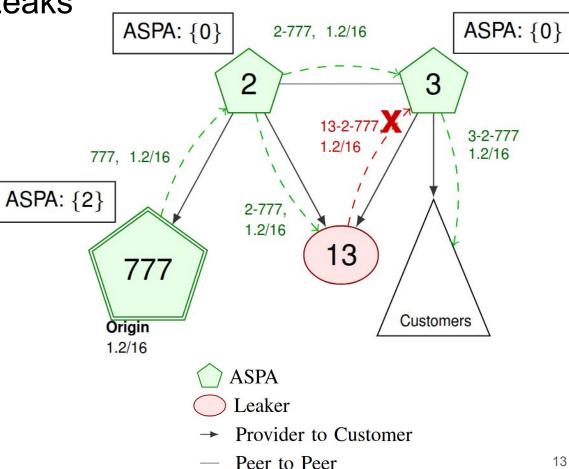
Enforce-first-as & First-ASN-Stripping Hijack

- ASPA states that the first ASN in the AS-Path must be the ASN of the neighbor
 - This just makes common sense to check that your neighbor is the last ASN on the AS-Path
 - This is using the enforce-first-as option on routers
- However we discovered this behavior is not the default on several routing vendors!
- What's worse, this is a global setting for Cisco
 - A bug for over 10 years, updated 2023
- Attackers can exploit this in what we call the first-asn-stripping hijack, where they strip their own ASN from the AS-Path
 - This reduces the AS-path length and attracts more traffic to the attacker

Vendor	Default	Global
Cisco	Y	Y
Juniper	N	N
Arista	Y	Ν
BIRD	N	N

ASPA Against Route Leaks

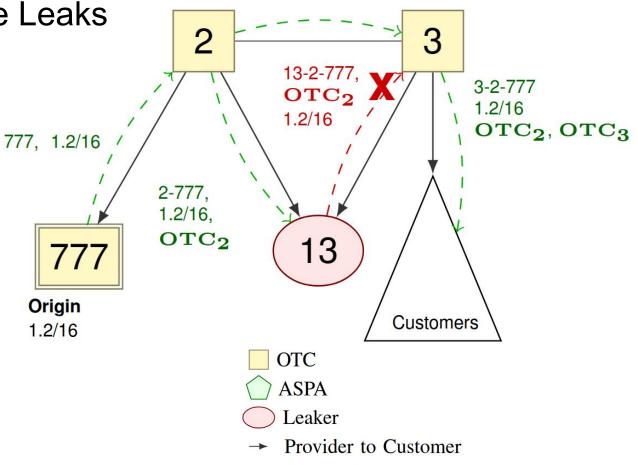
- Beyond hijacks, ASPA also protects against route leakage
- These are commonly misconfigurations that violate valley free routing
- ASPA uses special algorithms outlined in the draft by the IETF to detect route leaks
 - In short ensure valley free routing is taking place



2-777, 1.2/16, OTC₂

OTC Against Route Leaks

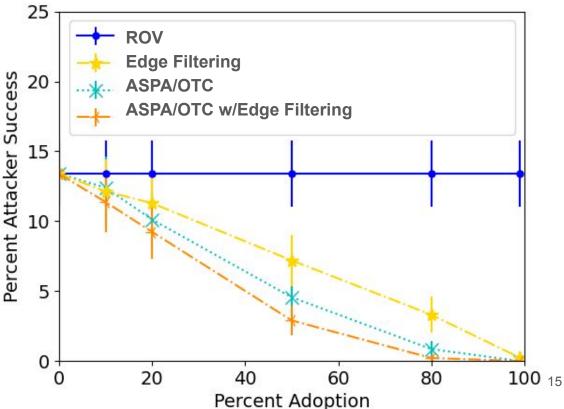
- OTC is another IETF proposed RFC
 - From a high level, 777 announcements with OTC should only be sent to customers
- OTC is already deployed, but only by a few smaller ASes



Peer to Peer

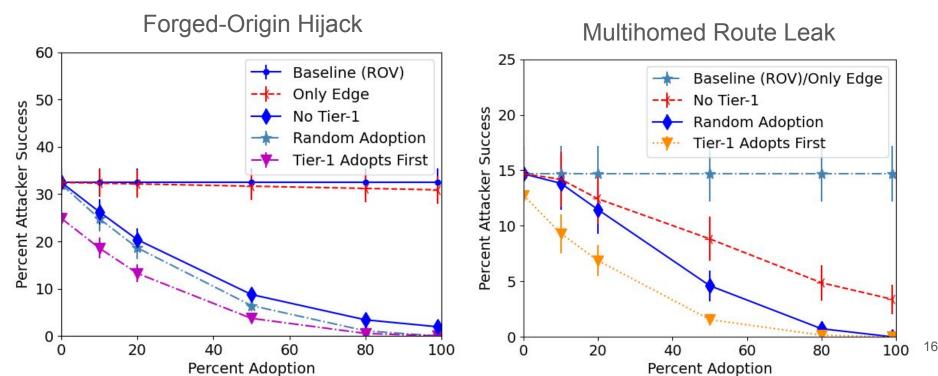
ASPA vs. OTC Route Leak Simulations

- We prove in a first that ASPA is equivalent to OTC for misconfigurations
- ASPA also protects against intentional leaks
 - (where an attacker would simply remove OTC)
- ASPA outperforms OTC when ASes drop transitive attributes
 - 1-3% of ASes do this



ASPA Adoption Scenarios

Contrary to prior works, we show that ASes benefit from ASPA even when tier-1 ASes do not adopt, motivating ASPA adoption throughout transit ASes and the internet at large



Reproducibility

All our simulations and everything in this work is **fully reproducible**

- Our simulations run off of BGPy, a leading Python BGP Security Simulator: <u>github.com/jfuruness/bgpy_pkg</u>
 - BGPy has been used in 5+ publications already with many more ongoing
 - It has been used by teams all around the world including NIST!
 - It is actively maintained by me, any questions just reach out :)
- We open-source our ROV collection tool here: <u>github.com/jfuruness/rov_collector</u>
 - We parse real-world ROV data from every known ROV data source
- We also open-source our code to run our simulations:
 - o github.com/jfuruness/aspa_eval
- We'd love to hear from you if you want to try it out, extend this, and implement some simulations to answer any unanswered questions!

Thank You!

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- Links:
 - BGPy:
 - github.com/jfuruness/bgpy_pkg
 - ROV Collector:
 - github.com/jfuruness/rov_collector
 - ASPA Eval:
 - github.com/jfuruness/aspa_eval
- I was also limited due to the time constraints, for more (especially comparison to other policies), read our paper :)
- Any questions?